

Introduction to quantum computing and simulability

Demonstrating a quantum advantage I

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Outline: Demonstrating a quantum advantage I

- Introduction and motivation;
- Linear optics;
 - Linear optics with bosons and fermions;
- BosonSampling;
- Random quantum circuits;

Review: postBPP and postBQP

• Postselection: The ability to condition acceptance on some particular (not-impossible) event, <u>no matter how unlikely</u>.



Quantum or (randomized) classical circuit

Review: postBPP and postBQP



Recipe for demonstrating quantum advantage

Take a restricted model of quantum computing A.
 e.g. circuits of commuting gates or linear optics

2 - Give it postselection, and see what comes out. (call it **postA**)

3 - If **A** + post-selection includes quantum computing, then **postA** = **postBQP**

4 - Suppose there is a classical algorithm to efficiently simulate A (i.e. sample from same distribution).
 Then postA ⊆ postBPP.

5 - But then **postBQP** ⊆ **postBPP** and **PH collapses!**

Quantum advantage (AKA "quantum supremacy")

- Show a quantum computer doing <u>something</u> a classical computer cannot.
 - Even if that "something" does not have any clear applications!
- **Task**: simulate some restricted quantum circuit;

i.e. nonuniversal(e.g. IQP, Linear optics)

Weak but exact simulation!

- Postselection argument: there is no efficient classical algorithm for this, otherwise PH collapses; ③
- Bad news: a realistic quantum device can't do it either! 🙁

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Linear optics: states

• System of *m* bosonic **modes** described by creation operators:

 $\{a_i^{\dagger}\}_{i=1,\ldots,m}$

• where

$$[a_i, a_j^{\dagger}] = \delta_{ij}$$

$$[a_i, a_j] = 0$$

$$a_i | \rangle = 0 \quad \text{Vacuum state}$$

• *n*-boson states:

$$|S\rangle = |s_1, s_2, \dots, s_m\rangle = \frac{(a_1^{\dagger})^{s_1} (a_2^{\dagger})^{s_2} \dots (a_m^{\dagger})^{s_m}}{\sqrt{s_1! s_2! \dots s_m!}}|\rangle$$

Linear optics: dynamics

• Linear-optical transformations (or interferometers):

$$U \in SU(m) \implies a_i \mapsto \sum_{j=1}^m U_{ij} a_j$$

- **Very** restricted set of operations!
 - The state space for *n* photons in *m* modes has dimension $\binom{m+n-1}{n}$
 - Also: not universal for QC!
- Alternative (Hamiltonian) characterization:

$$H = \sum_{ij} h_{ij} (a_i a_j^{\dagger} + a_j a_i^{\dagger})$$

Linear optics: dynamics

• Elementary two-mode linear-optical transformations:



Linear optics: dynamics

• Any *m*-mode interferometer can be written as a **circuit**:



Linear optics: example (HOM effect)

Consider a 50:50 (or balanced) beam splitter

• If one photons enters it...

Input: $|10\rangle = a_1^{\dagger}|\rangle$

Output:

$$\frac{1}{\sqrt{2}}(a_1^{\dagger} + ia_2^{\dagger}) |\rangle$$
$$= \frac{1}{\sqrt{2}}(|10\rangle + i|01\rangle)$$



Linear optics: example (HOM effect)

Consider a 50:50 (or balanced) beam splitter

• Now if two photons enter...

Input: $|11\rangle = a_1^{\dagger}a_2^{\dagger}|\rangle$

Output:
$$\frac{1}{2}(a_1^{\dagger} + ia_2^{\dagger})(ia_1^{\dagger} + a_2^{\dagger})|\rangle$$
$$= \frac{i}{\sqrt{2}}(|20\rangle + |02\rangle)$$



Hong-Ou-Mandel effect!

• If we had a general 2-mode transformation:

Two photons enter; what is the probability they leave separately?

Input: $|11\rangle = a_1^{\dagger}a_2^{\dagger}|\rangle$

Output:
$$(aa_1^{\dagger} + ba_2^{\dagger})(ca_1^{\dagger} + da_2^{\dagger})|\rangle$$

= $(ad + bc)|11\rangle + |\text{collision terms}\rangle$

• If we had a general 2-mode transformation:

• Probability of outcome $|11\rangle$:

$$\Pr(11) = |ad + bc|^2 = \left| \Pr\begin{pmatrix} a & b \\ c & d \end{pmatrix} \right|^2$$

$$\operatorname{Per}(A) = \sum_{\sigma \in S_m} \prod_{i=1}^m a_{i,\sigma(i)}$$

Side note: Fermionic "linear optics"

• What if we had **fermions**, rather than bosons?

$$U = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$$

Bosons
$$[a_i, a_j^{\dagger}] = \delta_{ij}$$

$$Per \begin{pmatrix} a & b \\ c & d \end{pmatrix}$$

Det $\begin{pmatrix} a & b \\ c & d \end{pmatrix} = 1$

Side note: Fermionic "linear optics"

• What if we had **fermions**, rather than bosons?

$$U = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & i \\ i & 1 \end{pmatrix}$$



• General n-photon m-mode interferometer U:



* from now on: no-collision inputs/outputs

• Transition between two (no-collision) states.



• Probability: $\Pr_{S \to T} = |\operatorname{Per}(U_{S,T})|^2$

 $U_{S,T}$: submatrix of U with rows/columns chosen according to S and T

Example:

Input: $|S\rangle = |110\rangle$

Output: $|T\rangle = |011\rangle$







$$\Pr_{S \to T} = |dh + ge|^2$$



Side note: Fermionic linear optics II

Transition:
$$|S\rangle \rightarrow |T\rangle$$

Bosons: $\Pr_{S \rightarrow T} = |\operatorname{Per}(U_{S,T})|^2$
Fermions: $\Pr_{S \rightarrow T} = |\operatorname{Det}(U_{S,T})|^2$
Permanent and determinant are similar.
but in terms of complexity, they are
very different!
EXP
|
PSPACE
Permanent
PP = postBQP
PH
|
postBPP
BQP
NP
L
Determinant

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It is not hard to show that

Linear optics + postselection = **postBQP**



It is not hard to show that

Linear optics + postselection = **postBQP**

 \Rightarrow Linear optics cannot be simulated on a classical computer!

 But this is about exact simulation - a real-world experiment can't simulate this exactly either!

Time to bring out the (complexity-theoretic) big guns!



• Consider the following *n*-photon *m*-mode experiment:



where $m = O(n^2)$ and U is some uniformly-random matrix.

- BosonSampling task: sample from that n-photon distribution;
 - Or in fact any sufficiently close to it!
 - **Not** a decision problem!

Theorem [Aaronson and Arkhipov]

If there was a classical algorithm capable of sampling efficiently from some distribution D' such that

 $|D - D'| < \delta$

in time $poly(n, 1/\delta)$, the polynomial hierarchy (**PH**) would collapse to its 3rd level!

- Outline of result:
 - 1. Probabilities = permanents of submatrices of a random matrix;



- Outline of result:
 - 1. Probabilities = permanents of submatrices of a random matrix;
 - 2. Permanent-of-Gaussian conjecture;

Fact: Permanent is **#P**-hard in the worst case.

<u>Conjecture</u>: Gaussian matrices are among the hardest Permanents to compute;

Plausible conjecture that has **nothing** to do with linear optics!

- Outline of result:
 - 1. Probabilities = permanents of submatrices of a random matrix;

2. Permanent-of-Gaussian conjecture;

3. A classical simulation that samples from a distribution very close to the ideal one must get many of these **#P-hard** probabilities right!

. . . .

- 4. Another minor conjecture + 90 pages of complexity-theoretic cannons shooting at the problem...
- 5. If an efficient classical simulation for approximate BosonSampling existed, **PH** would collapse to the 3rd level!

- Some additional comments:
 - Result does **not** mean linear optics can compute permanents!
 - $m=O(n^2)$ means no-collision outcomes vastly more likely.
 - Randomness of *U* avoids structures that classical algorithms could "exploit".
 - This sampling task has no known applications;
 Except as a "demonstration of force" by quantum devices ^(C)
 - Best known classical algorithm takes time $O(n 2^n)$;

BosonSampling: pros and cons

- Cons
 - No error-correction, so it is unclear whether experiments could be performed to this level of accuracy!
 - Hard to verify that device is doing what it should;
 - Leandro has more to say about this!
- Pros
 - Much "easier" to implement than universal quantum computation;
 - Doesn't require nonlinearities or adaptive measurements;
 - ~ 50-90 photon experiment?
 - New insights into foundations of q. computing and q. optics;

First experiments

• Four small-scale experiments reported in December 2012.

Brisbane

Rome

Science

Photonic Boson Sampling in a Tunable Circuit

Matthew A. Broome,^{1,2}* Alessandro Fedrizzi,^{1,2} Saleh Rahimi-Keshari,² Justin Dove,³ Scott Aaronson,³ Timothy C. Ralph,² Andrew G. White^{1,2}

nature photonics

ETTERS

PUBLISHED ONLINE: 26 MAY 2013 | DOI: 10.1038/NPHOTON.2013.112

Integrated multimode interferometers with arbitrary designs for photonic boson sampling

Andrea Crespi^{1,2}, Roberto Osellame^{1,2}*, Roberta Ramponi^{1,2}, Daniel J. Brod³, Ernesto F. Galvão³*, Nicolò Spagnolo⁴, Chiara Vitelli^{4,5}, Enrico Maiorino⁴, Paolo Mataloni⁴ and Fabio Sciarrino⁴*

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Boson Sampling on a Photonic Chip

Justin B. Spring,¹* Benjamin J. Metcalf,¹ Peter C. Humphreys,¹ W. Steven Kolthammer,¹ Xian-Min Jin,^{1,2} Marco Barbieri,¹ Animesh Datta,¹ Nicholas Thomas-Peter,¹ Nathan K. Langford,^{1,3} Dmytro Kundys,⁴ James C. Gates,⁴ Brian J. Smith,¹ Peter G. R. Smith,⁴ Ian A. Walmsley¹*



Experimental boson sampling

PUBLISHED ONLINE: 12 MAY 2013 | DOI: 10.1038/NPHOTON.2013.102

Max Tillmann^{1,2}*, Borivoje Dakić¹, René Heilmann³, Stefan Nolte³, Alexander Szameit³ and Philip Walther^{1,2}*

First experiments

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and Philip Walther^{1,2*}

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A zoo of intermediate quantum models

- Many other similar "quantum supremacy" results!
 - IQP; [BJS]
 - Constant-depth quantum circuits [TD]; BosonSampling
 - Quantum approximate optimization algorihtms; [Farhi, Harrow]
 - A version of the 1-clean-qubit model; [Morimae, Fujii, Fitzsimmons]

These predate

- Generalized Clifford circuits [JVdN]
- Random quantum circuits;
- And more!

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Random quantum circuits

- Google's Quantum Al Lab + John Martini's group at UCSB
 - Another way to demonstrate the power of quantum computers.
- Meet Google's 72-qubit quantum processor: Bristlecone!



• Goal: Demonstrate "Quantum advantage" on 72 qubits! (...soon)

Image: Google AI blog



- But what is known about the complexity of simulating this?
- Aaronson and Chen [Dec '16]
 - Task: "Given as input a random quantum circuit C, generate a set of outputs such that at least 2/3 of them have above-median probability in the output distribution of C";



Generate a list of samples where $\geq 2/3$ of them are among those above this height;

- But what is known about the complexity of simulating this?
- Aaronson and Chen [Dec '16]
 - Task: "Given as input a random quantum circuit C, generate a set of outputs such that at least 2/3 of them have above-median probability in the output distribution of C";
 - Needs one further conjecture (similar to permanent-of-Gaussians in BosonSampling);
- Bouland, Fefferman, Nirkhe and Vazirani [March '18!]
 - Made a major step towards improving the status of the required conjectures.

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- Aaronson and Chen [Dec '16]
 - Neat ideal Let's build one already, what's Task: "Given as input a rar set of obability outputs such the in the to permanent-of-Gaussians in Nee Bosd
- merman, Nirkhe and Vazirani [March '18!] Bouland
 - Made a major step towards improving the status of the required conjectures.